Improvements to CHvote

Towards an end-to-end verifiable voting system



Département de la sécurité et de l'économie Direction générale des systèmes d'information

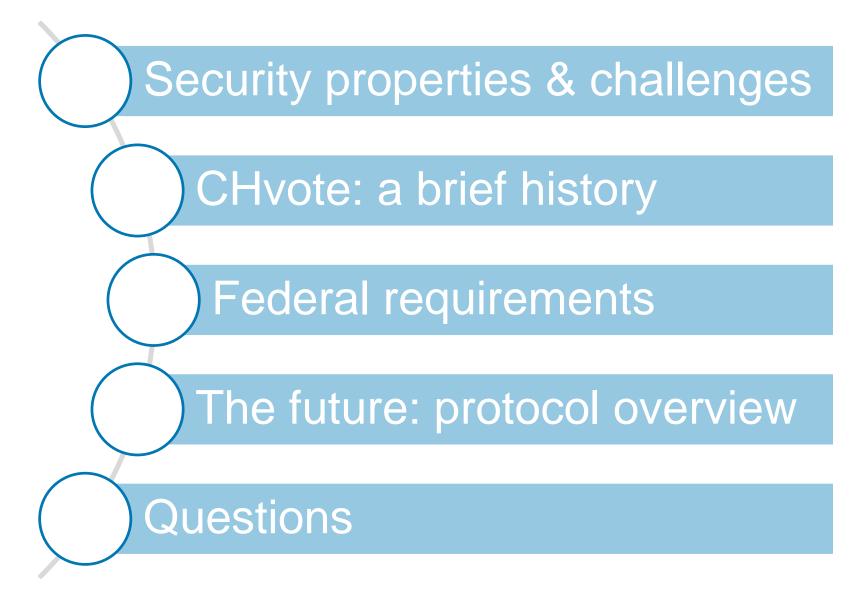


Short Bio

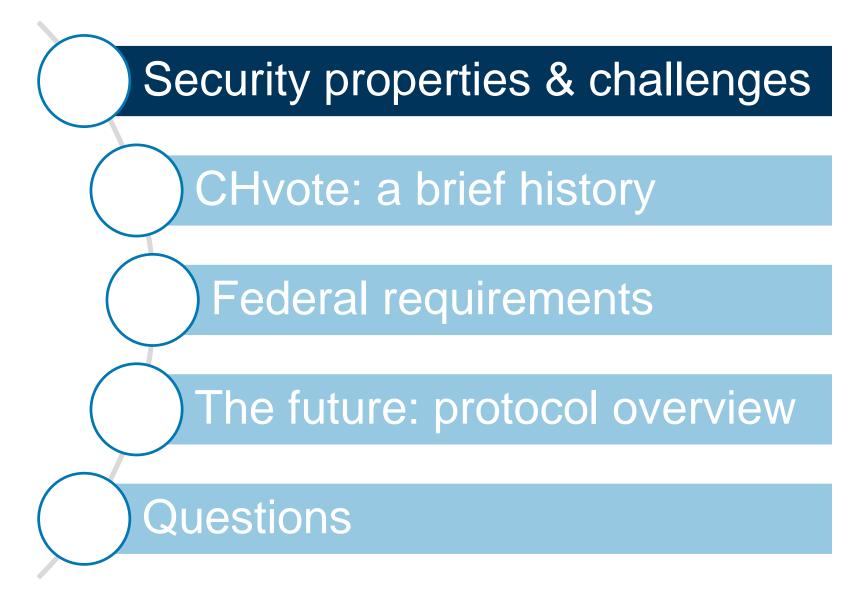
Thomas Hofer / @thhofer / thomas.hofer@etat.ge.ch

- EPFL MSc in IT
- IT / Java consultant
- Now
 - → Internet voting cryptography @ State of Geneva
 - → Java DEV & AppSec
- Outside from work
 - → OWASP-Geneva co-chapter leader
 - → Married, 2 kids

Outline



Outline



Security properties

Target security properties





Availability



Voter authentication



No early tally

Security challenges

Partially contradicting requirements and other challenges

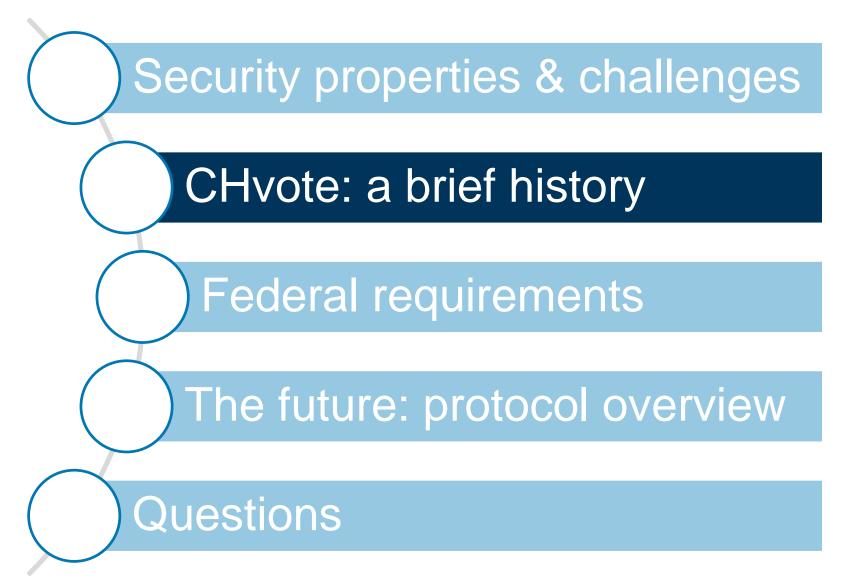
- Vote secrecy vs. result integrity
 - → Cryptographically challenging (but feasible)
- Enfranchisement vs. authentication
 - → Typically opposed
 - → But: in CH, voting legitimation cards are sent to voters (Swiss Post is trusted)
 - → For mail-in ballots / polling station voting:
 - Voting card + signature + DOB
 - → For internet voting:
 - Secrets printed on voting card + DOB

Security challenges

Partially contradicting requirements and other challenges (ctd.)

- Availability
 - → OK, but… DDOS??
 - → Standard technical counter-measures
 - → Internet voting closes 24 hours before polling stations





The past of CHvote

First generation E-Voting system

- 2001: start of project
- 2003: first use



• Partners



The present of CHvote

Individual verifiability & major appearance overhaul



financement et de l'aménagement de l'infrastructure ferroviaire (Contreprojet direct à l'initiative populaire "Pour les transports publics", qui a été retirée) ?

Z8H5

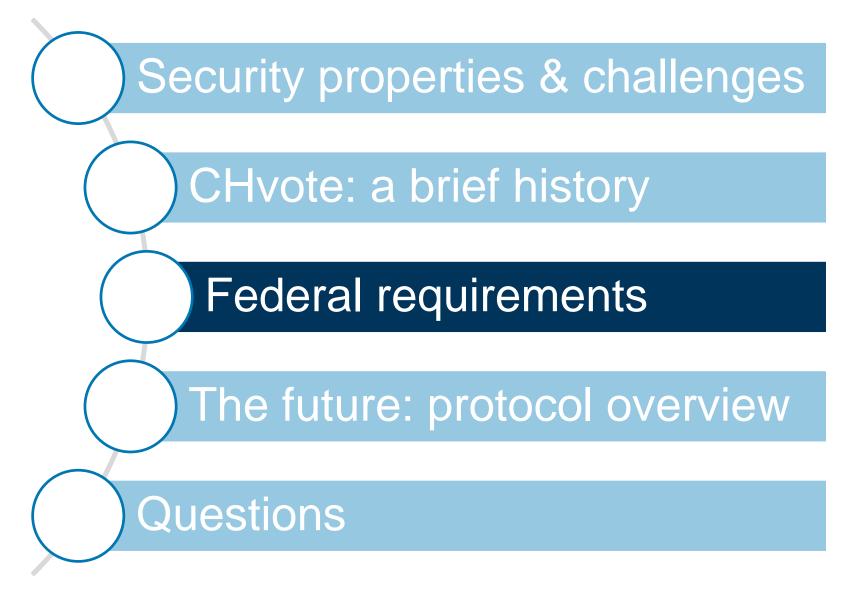
J5B9

The future of CHvote

End-to-end verifiable internet voting protocol

- New academic partnerships
 - → Berner Fachhochschule
 - → INRIA / Bristol
 - → ITU Copenhagen
 - → EPFL
- New cryptographic protocol
 - → End-to-end encryption
 - → Universal Verifiability
 - → Control Components
- Currently in development, ETA: 2019





New Ordinance on Electronic Voting

- Published in 2013, enacted 2014
 - → Collaborative work between lawmakers, academia and operating staff
- Compliance levels
 - → The higher the compliance, the more voters allowed
- Reference

https://www.bk.admin.ch/themen/pore/evoting/07979/index.html

Individual Verifiability

Voters must receive proof that the server system has registered the vote as it was entered by the voter on the user platform – *VEleS, art. 4*

Liste de codes pour la carte n° 5874-8863-1400-8743		•	
Votation fédérale			
Question 1 Acceptez-vous l'arrêté fédéral du 20 juin 2013 portant règlement du financement et de l'aménagement de l'infrastructure ferroviaire (Contre- projet direct à l'initiative populaire "Pour les transports publics", qui a été retirée) ?	Oui A2B4	Non J5B9	Blanc Z8H5

End-to-End Encryption

Votes must not be stored or transmitted in unencrypted form at any time from being entered to tallying. – *Technical and administrative requirements, section* 3.3.4

Universal Verifiability

For universal verification, the auditors receive proof that the result has been ascertained correctly. They must evaluate the proof in a observable procedure. – *VEleS, art. 5 paragraph 4*

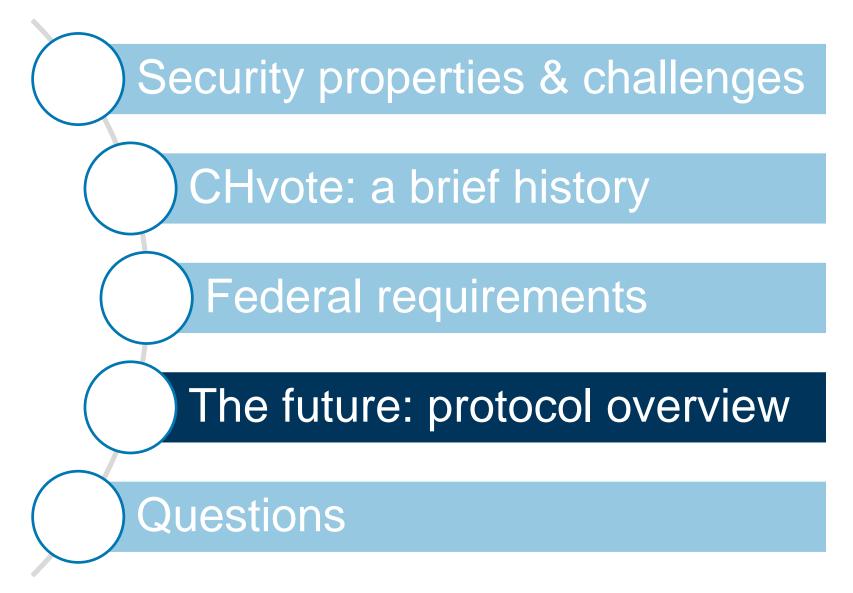
Control Components

The trustworthy part of the system includes either one or a small number of groups of independent components secured by special measures (control components). Their use must also make any abuse recognisable if per group only one of the control components works correctly and in particular is not manipulated unnoticed. – *VEleS, art. 5, par. 6*

Compliance levels

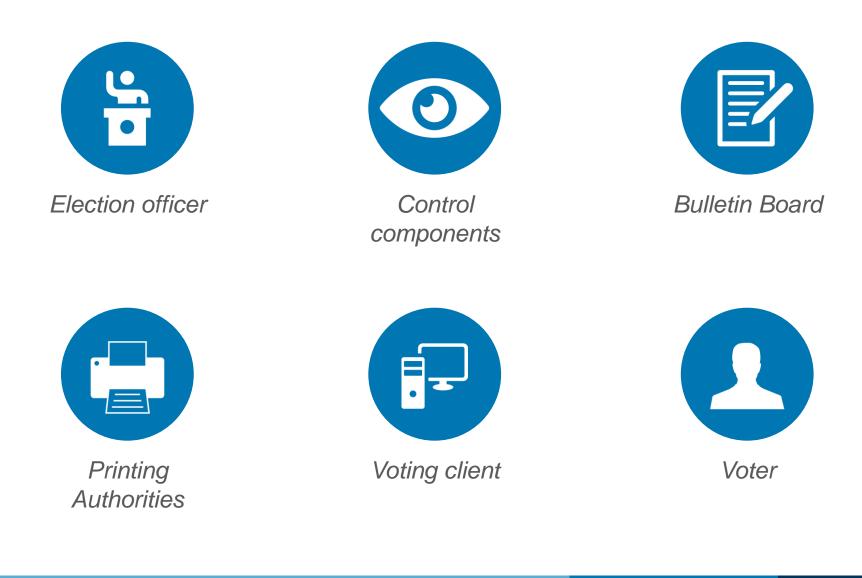
- First level
 - → Individual verifiability
 - → Internet voting for up to 30% of voters
- Second level
 - → Add certifying audit
 - → Internet voting for up to 50% of voters
- Third level
 - Add universal verifiability, control components and end-to-end encryption
 - → New certifying audit
 - → Internet voting for up to 100% of voters





Protocol actors

Stakeholders from the perspective of the cryptographic protocol



Key cryptographic primitives

A brief overview

- El Gamal homomorphic encryption
- Oblivious Transfer for individual verifiability
 - → <u>Cast-as-Intended Verification in Electronic Elections Based on</u> <u>Oblivious Transfer</u>
- Pedersen Commitments
- Non-interactive Zero-Knowledge Proofs (ZKP)
- Wikström's Proof of a Shuffle

Homomorphic encryption

What is it?

- Principles
 - → Operations performed on cipher texts
 - → Result visible on recovered plain texts
 - → Example:
 - Encrypt 2
 - Multiply cipher text by 3
 - Decrypt
 - Result is 6
- For this project: El Gamal encryption

Homomorphic encryption

How and why?

- Used for voter credentials
 - → Voter authentication
- Used for encrypting the ballots
 → Vote secrecy
- Allows re-encryptions
 - → Useful for anonymizing when shuffling
 - → Vote secrecy
- Allows for key sharing
 - → Control components each hold a key share
 - → Vote secrecy & result integrity

Oblivious Transfer

What does it mean and why is it useful?

- In short
 - → Server knows n secret messages
 - → Client allowed to retrieve k secret messages
 - → Server cannot know which messages the client asked for
 - → Perfect match for the verification codes issue!
 - → Vote secrecy & Result integrity
- In detail
 - → <u>Cast-as-Intended Verification in Electronic Elections Based on</u> <u>Oblivious Transfer</u>

Commitments and ZKPs

How and why?

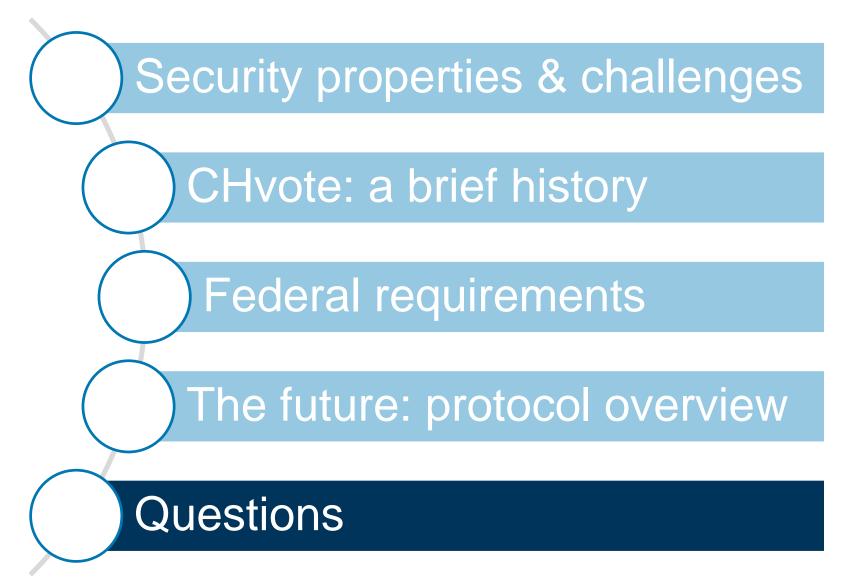
- "public" commitments for the secrets
 - → Share a value computed from secret, without leaking info
- ZKPs relative to those commitments
 - → Prove that
 - Secret value used in computation =
 - secret value used for commitment
 - → Chain of truth from key generation to ballot decryption
- Combination yields Universal verifiability
 - → Result integrity

Wikström's Proof of a Shuffle

Why?

- Re-encrypting mix-net
 - → Each component re-encrypts each ballot and shuffles them
- Since shuffled, simple pre-image proofs would not work
- Since re-encrypted, ciphertexts are not equal
 → Vote secrecy
- Need for a specific proof that the cryptographic shuffle is valid
 - → Result integrity





Further reading

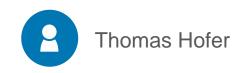
And references

- Published protocol specification
 - → <u>https://eprint.iacr.org/2017/325</u>
- Published PoC code
 - https://github.com/republique-et-canton-de-geneve/chvoteprotocol-poc
- Federal requirements

<u>https://www.bk.admin.ch/themen/pore/evoting/07979/index.html</u>



Thank you!





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